

# Bounded Genericity

Ludwig Staiger

Lehrstuhl für Theoretische Informatik  
Institut für Informatik  
Martin-Luther-Universität Halle-Wittenberg

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- 2 Structure Functions of Languages
- 3 Dimensions of  $\omega$ -languages

## Notation: Strings and Languages

- Finite Alphabet  $X = \{0, \dots, r - 1\}$ , cardinality  $|X| = r \geq 2$
- Finite strings (words):  $w = x_1 \cdots x_l \in X^*$ ,  $x_i \in X$ , length  $|w| = l$
- Languages:  $W \subseteq X^*$
- Infinite strings ( $\omega$ -words):  $\xi = x_1 \cdots x_l \cdots \in X^\omega$
- Prefixes of infinite strings:  $\xi[0..n] \in X^*$ ,  $|\xi[0..n]| = n$
- Prefix relation:  $w \sqsubseteq \eta$ , for  $\eta \in X^* \cup X^\omega$
- Language of prefixes:  $\mathbf{pref}(B) := \{w : \exists \eta (\eta \in B \wedge w \sqsubseteq \eta)\}$

# $X^\omega$ as CANTOR space

**Metric:**  $\rho(\eta, \xi) := \inf \{r^{-|w|} : w \sqsubset \eta \wedge w \sqsubset \xi\}$

**Balls in  $(X^\omega, \rho)$ :**  $w \cdot X^\omega = \{\eta : w \sqsubset \eta\}$

**Diameter:**  $\text{diam } w \cdot X^\omega = r^{-|w|}$

**Open sets:**  $W \cdot X^\omega = \bigcup_{w \in W} w \cdot X^\omega$

**Closure:**  $\mathcal{C}(F) = \{\xi : \mathbf{pref}(\xi) \subseteq \mathbf{pref}(F)\}$

## Property

$F \subseteq X^\omega$  is closed if and only if  $\mathbf{pref}(\xi) \subseteq \mathbf{pref}(F)$  implies  $\xi \in F$ .

# Extension functions (Ambos-Spies & Busse '03)

**Extension function:**  $f : X^* \rightarrow X^*$

**$k$ -bounded extension function:**  $f : X^* \rightarrow X^k$

$f : X^* \rightarrow X^*$  **meets**  $\xi \in X^\omega$ :  $\exists w (w \sqsubset \xi \wedge w \cdot f(w) \sqsubset \xi)$

$$M_f := \bigcup_{w \in X^*} w \cdot f(w) \cdot X^\omega,$$

$$F_f := X^\omega \setminus M_f$$

$$M_{\mathcal{F}} := \bigcap_{f \in \mathcal{F}} M_f,$$

$$F_{\mathcal{F}} := X^\omega \setminus M_{\mathcal{F}} = \bigcup_{f \in \mathcal{F}} F_f$$

$\xi$  is  $(f_n)_{n \in \mathbb{N}}$ -**generic**:  $\xi \in \bigcap_{n \in \mathbb{N}} M_{f_n} = X^\omega \setminus \bigcup_{n \in \mathbb{N}} F_{f_n}$

## $k$ -bounded automatic genericity

$f : X^* \rightarrow X^k$  is defined by a finite automaton  $\mathcal{A} = (X, Q, q_0, \delta)$  and a function  $g : Q \rightarrow X^k$  via

$$f(w) := g(\delta(q_0, w)).$$

### Lemma

*If  $f$  is a  $k$ -bounded automatic extension function then  $F_f$  is definable by a finite automaton.*

### Theorem (Ambos-Spies & Busse '03)

*Let  $\{f_n^k : n \in \mathbb{N}\}$  be the family of all  $k$ -bounded automatic extension functions. Then*

$$\bigcup_{n \in \mathbb{N}} F_{f_n^k} \subset \bigcup_{n \in \mathbb{N}} F_{f_n^{k+1}}.$$

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## Special case: Forbidden subwords

For some  $v \in X^*$ ,  $f \equiv v$  is a constant function, that is,

$$E(\neg v) := F_f = X^\omega \setminus X^* \cdot v \cdot X^\omega.$$

### Definition (Overlap-free words)

A word  $u \in X^*$  is **overlap-free** if and only if no non-empty proper prefix of  $u$  is a suffix of  $u$ .



## A hierarchy result by VOLKMANN

### Theorem (Volkmann '53)

1. **length  $k$ :** Let  $v \in X^k$  be overlap-free, and  $u \in X^k$  neither overlap-free nor of the form  $a^k$ ,  $a \in X$ . Then

$$\dim_{\mathbb{H}} E(\neg v) < \dim_{\mathbb{H}} E(\neg u) < \dim_{\mathbb{H}} E(\neg 0^k).$$

2. **length  $k$  and  $k + 1$ :** Let  $u \in X^{k+1}$ . If  $|X| \geq 3$  or  $u$  is not overlap-free then

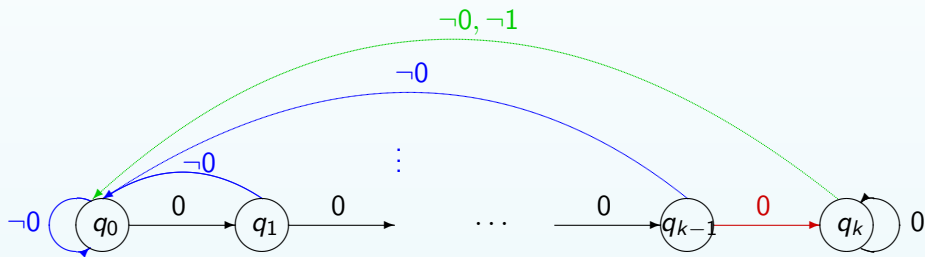
$$\dim_{\mathbb{H}} \bigcup_{|v|=k} E(\neg v) < \dim_{\mathbb{H}} E(\neg u).$$

# The rôle of the alphabet

**Observation.** For  $X = \{0, 1\}$  we have

$$\dim_{\mathbb{H}} E(\neg \mathbf{0}^k \mathbf{1}) = \dim_{\mathbb{H}} E(\neg \mathbf{0}^k)$$

$$E(\neg \mathbf{0}^k \mathbf{1}) \subseteq \bigcup_{n \in \mathbb{N}} F_{f_n^k} \quad (f_n^k \text{ is } k\text{-bounded automatic})$$



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# Structure function

Definition (Structure Function of  $W \subseteq X^*$  and  $F \subseteq X^\omega$ )

$$s_W(n) := |\{w : w \in W \wedge |w| = n\}| = |W \cap X^n|$$

$$s_F := s_{\text{pref}(F)}$$

## Special case: $E(\neg 0^{k-1}1)$ and $E(\neg 0^k)$

$$\begin{aligned} \underline{E}_k &:= E(\neg 0^{k-1}1) = X^\omega \setminus X^* \cdot \mathbf{0}^{k-1}\mathbf{1} \cdot X^\omega \\ E_k &:= E(\neg 0^k) = X^\omega \setminus X^* \cdot \mathbf{0}^k \cdot X^\omega \end{aligned}$$

Their structure functions satisfy

$$\begin{aligned} \underline{s}_k(n) &= s_k(n) = r^n, & \text{if } 0 \leq n < k \text{ and} \\ \underline{s}_k(k) &= s_k(k) = r^k - 1, & \end{aligned} \quad (1)$$

and the recurrences

$$\begin{aligned} \underline{s}_k(n) &= r \cdot \underline{s}_k(n-1) - \underline{s}_k(n-k), & \text{if } n > k \\ s_k(n) &= (r-1) \cdot \sum_{i=1}^k s_k(n-i), & \text{if } n > k. \end{aligned} \quad (2)$$

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## Asymptotic behaviour of $\underline{s}_k$ and $s_k$

Let  $\underline{\lambda}_k$  and  $\lambda_k$  be the maximum roots of the polynomials

$$\begin{aligned} \underline{p}_k(z) &= z^k - r \cdot z^{k-1} + 1 \text{ and} \\ p_k(z) &= z^k - \sum_{i=0}^{k-1} (r-1) \cdot z^i, \text{ respectively.} \end{aligned}$$

**Remark.**  $\underline{\lambda}_k < \lambda_k$ , and if  $r > 2$  then  $\lambda_k < \lambda_{k+1}$ .

$$\underline{s}_k(n) \approx \underline{\lambda}_k^n \quad \text{and} \quad s_k(n) \approx \lambda_k^n$$

# Overlapping

## Definition

Let  $f : X^* \rightarrow X^k$  be an extension function.

- ① We call  $f$  **overlap-free** :  $\iff$   
 $\forall w \forall v (v \sqsubset w \sqsubset v \cdot f(v) \rightarrow v \cdot f(v) \not\sqsubseteq w \cdot f(w)).$
- ② We call  $f$  **fully overlapping** :  $\iff$   
 $\forall w \forall v (v \sqsubset w \sqsubset v \cdot f(v) \rightarrow v \cdot f(v) \sqsubseteq w \cdot f(w)).$

## Lemma

Let  $f : X^* \rightarrow X^k$  be an extension function.

- ① *There is an overlap-free extension function  $g : X^* \rightarrow X^k$  such that  $F_g \subseteq F_f$ .*
- ② *There is a fully overlapping extension function  $h : X^* \rightarrow X^k$  such that  $F_h \supseteq F_f$ .*



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## Minimal and maximal elements

$$\mathcal{E}_k := \{F_f : f : X^* \rightarrow X^k\}$$

### Lemma

If  $F_f$  is a minimal (maximal) w.r.t. set inclusion element in  $\mathcal{E}_k$  then there is an overlap-free extension function  $g : X^* \rightarrow X^k$  (a fully overlapping extension function  $h : X^* \rightarrow X^k$ ) such that  $F_f = F_g$  ( $F_f = F_h$ ).

### Theorem

- 1 If  $f : X^* \rightarrow X^k$  then  $\underline{s}_k \leq s_{F_f} \leq s_k$ .
- 2 If  $g : X^* \rightarrow X^k$  is overlap-free then  $s_{F_g} = \underline{s}_k$  and  $F_g$  is minimal in  $\mathcal{E}_k$ .
- 3 If  $h : X^* \rightarrow X^k$  is fully overlapping then  $s_{F_h} = s_k$  and  $F_h$  is maximal in  $\mathcal{E}_k$ .

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# Minkowski or box-counting dimensions

$$\underline{\dim}_{\mathbf{B}} F := \liminf_{n \rightarrow \infty} \frac{1}{n} \cdot \log_{|\mathcal{X}|}(s_F(n) + 1)$$

$$\overline{\dim}_{\mathbf{B}} F := \limsup_{n \rightarrow \infty} \frac{1}{n} \cdot \log_{|\mathcal{X}|}(s_F(n) + 1)$$

## Property

$$\underline{\dim}_{\mathbf{B}} F = \underline{\dim}_{\mathbf{B}} \mathcal{C}(F)$$

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$$E \subseteq F \rightarrow \underline{\dim}_{\mathbf{B}} E \leq \underline{\dim}_{\mathbf{B}} F$$

$$\overline{\dim}_{\mathbf{B}} E \cup F = \max \{ \overline{\dim}_{\mathbf{B}} E, \overline{\dim}_{\mathbf{B}} F \}$$

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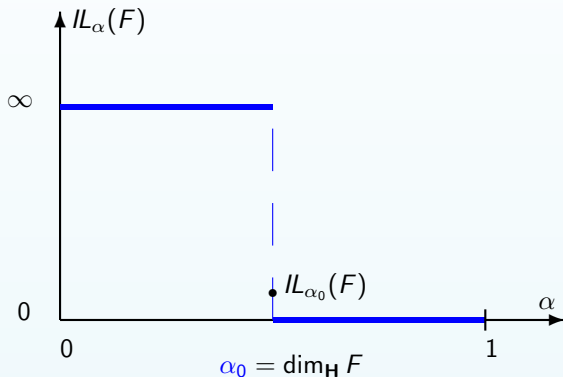
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# Hausdorff measure

$$IL_\alpha(F) := \lim_{l \rightarrow \infty} \inf \left\{ \sum_{w \in W} r^{-\alpha \cdot |w|} : F \subseteq \bigcup_{w \in W} w \cdot X^\omega \wedge \forall w (w \in W \rightarrow |w| \geq l) \right\}$$



# Hausdorff dimension

$$\dim_{\mathbf{H}} F = \sup \{ \alpha : IL_{\alpha}(F) = \infty \} = \inf \{ \alpha : IL_{\alpha}(F) = 0 \}$$

## Proposition

- 1  $\dim_{\mathbf{H}} \bigcup_{i \in \mathbb{N}} F_i = \sup \{ \dim_{\mathbf{H}} F_i : i \in \mathbb{N} \}$
- 2  $\dim_{\mathbf{H}} w \cdot F = \dim_{\mathbf{H}} F$
- 3  $\dim_{\mathbf{H}} F \leq \underline{\dim}_{\mathbf{B}} F \leq \overline{\dim}_{\mathbf{B}} F$



## Relations between the dimensions

Definition ("Subtree" rooted at  $w \in X^*$ )

$$F/w := \{\xi : w \cdot \xi \in F\}$$

Theorem (St.'89)

Let  $\log_r \gamma(n) = o(n)$  for some  $\gamma : \mathbb{N} \rightarrow \mathbb{N}$ . If  
 $s_{F/w}(n) \leq \gamma(|w|) \cdot s_F(n)$  for all  $w \in X^*$  and all  $n \in \mathbb{N}$  then

$$\dim_{\mathbb{H}} F = \overline{\dim}_{\mathbb{B}} F.$$

Observation

If  $f : X^* \rightarrow X^k$  is overlapfree or fully overlapping then  
 $s_{F_f/w}(n) \leq s_{F_f}(n)$ , for all  $w \in X^*$  and all  $n \in \mathbb{N}$ .

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## The dimensions of $F_f$

### Corollary

If  $f : X^* \rightarrow X^k$  is overlapfree or fully overlapping then

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Lemma (Mauldin & Williams '88, Bandt '89, St.'89)

If  $f : X^* \rightarrow X^k$  is automatic then

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## General hierarchy result I

### Theorem

*If  $f$  is a  $k$ -bounded extension function and  $f'$  is a  $(k + 1)$ -bounded extension function over  $X$  then  $F_f \not\subseteq F_{f'}$  and*

$$\log_{|X|} \underline{\lambda}_k \leq \dim_H F_f \leq \log_{|X|} \lambda_k \leq$$

$$\log_{|X|} \underline{\lambda}_{k+1} \leq \dim_H F_{f'} \leq \log_{|X|} \lambda_{k+1}.$$

*Equality of Hausdorff dimension  $\dim_H F_f = \dim_H F_{f'}$  can hold only if  $|X| = 2$ ,  $\dim_H F_f = \log_2 \lambda_k$  and  $\dim_H F_{f'} = \log_2 \underline{\lambda}_{k+1}$ .*

## General hierarchy result II

### Theorem

Let, for  $k \in \mathbb{N}$ ,  $\{f_n^k : n \in \mathbb{N}\}$  be a countable family of  $k$ -bounded extension functions containing all constant functions. Then

$$\bigcup_{n \in \mathbb{N}} F_{f_n^k} \not\subseteq \bigcup_{n \in \mathbb{N}} F_{f_n^{k+1}}.$$

### Corollary

Let  $\{f_n^k : n \in \mathbb{N}\}$  be the family of all  $k$ -bounded computable (recursive) extension functions. Then

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