

# Insertion in normal numbers

Verónica Becher

Universidad de Buenos Aires

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# Borel normal numbers

Let  $b$  be an integer greater than or equal to 2.

A real number is **normal to base  $b$**  if in its base- $b$  expansion every block of digits occurs with the same limiting frequency as every other block of the same length.

# Borel normal numbers

Let  $b$  be an integer greater than or equal to 2.

A real number is **normal to base  $b$**  if in its base- $b$  expansion every block of digits occurs with the same limiting frequency as every other block of the same length.

Instead of expansions of numbers we talk about **sequences of digits/symbols** that are normal to a  $b$ -alphabet.

# Modifying normal sequences

- ▶ **Selection of subsequences**  
Wall 1949, Agafonov 1968, Kamae and Weiss 1975
- ▶ **Sums**  
Rauzy 1976; Volkonoff 1979, Aistleitner 2017
- ▶ **Transformations by finite state transducers**  
Carton and Orduna 2020
- ▶ **Insertion in positions in a set of density zero**  
Figueira 2002, Aistleitner 2017

# Modifying effective Hausdorff dimension of sequences

The effective Hausdorff dimension of an infinite sequence  $x$  is a real number between 0 and 1 which measures the asymptotic information density of  $x$  (Lutz 2000, Mayordomo 2002).

Given a sequence  $x$  of effective Hausdorff dimension  $s$ , how much do we need to change  $x$  to obtain some  $y$  of dimension  $t$ ?

For binary sequences,

Greenberg, J.S.Miller, Shen and Westrick, 2018

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How can we change dimension but from binary sequences to ternary ?

# Today's question

## Enlarge the alphabet

Given a normal sequence, how can we **insert** symbols so that the expanded sequence is normal to the enlarged alphabet?

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Given a normal sequence, how can we **insert** symbols so that the expanded sequence is normal to the enlarged alphabet?

In other words, given a sequence of finite-state dimension 1, how can we **insert** symbols to obtain dimension 1 in the enlarged alphabet?



# Two theorems

## First theorem

Any-symbol insertion in some constructed normal sequences.

## Second theorem (Zylber 2017,2021)

Just the new symbol insertion in arbitrary normal sequences.

## Example of insertion

Theorem (Champernowne 1933)

*The concatenation of all blocks in length-lexicographic order is normal*

01 00 01 10 11 000 001 010 011 100 101 110 111 0000 0001 ...

# Example of insertion

## Theorem (Champernowne 1933)

*The concatenation of all blocks in length-lexicographic order is normal*

01 00 01 10 11 000 001 010 011 100 101 110 111 0000 0001 ...

If we enlarge the alphabet with one greater symbol,

01<sub>▲</sub>00 01<sub>▲</sub>10 11 000 001<sub>▲</sub>010 011<sub>▲</sub>100 101<sub>▲</sub>110 111 0000 0001 ...

01<sub>2</sub> 00 01 0<sub>2</sub> 10 11 1<sub>2</sub> 20 21 2<sub>2</sub> 000 001 00<sub>2</sub> 010 011 01<sub>2</sub> 020 021 02<sub>2</sub> 100 101 10<sub>2</sub> 110  
111 11<sub>2</sub> 120 121 12<sub>2</sub> 200 201 20<sub>2</sub> 210 211 21<sub>2</sub> 220 221 22<sub>2</sub> 0000 0001...

## Observation

Consider concatenation in lexicographic order of all blocks of length  $n$ . Viewed circularly, each block of length  $n$  occurs exactly  $n$  times at positions different modulo  $n$ .

	positions	
$n=2$	12 34 56 78	
	00 01 10 11	
	00 01 10 11	00 occurs twice, at positions different modulo 2

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	positions	
n=2	12 34 56 78	
	00 01 10 11	
	00 01 10 11	00 occurs twice, at positions different modulo 2
	00 01 10 11	
	00 01 10 11	01 occurs twice, at positions different modulo 2

## Observation

Consider concatenation in lexicographic order of all blocks of length  $n$ . Viewed circularly, each block of length  $n$  occurs exactly  $n$  times at positions different modulo  $n$ .

	positions	
n=2	12 34 56 78	
	00 01 10 11	
	00 01 10 11	00 occurs twice, at positions different modulo 2
	00 01 10 11	
	00 01 10 11	01 occurs twice, at positions different modulo 2
	00 01 10 11	
	00 01 10 11	10 occurs twice, at positions different modulo 2

## Observation

Consider concatenation in lexicographic order of all blocks of length  $n$ . Viewed circularly, each block of length  $n$  occurs exactly  $n$  times at positions different modulo  $n$ .

	positions	
n=2	12 34 56 78	
	00 01 10 11	
	00 01 10 11	00 occurs twice, at positions different modulo 2
	00 01 10 11	
	00 01 10 11	01 occurs twice, at positions different modulo 2
	00 01 10 11	
	00 01 10 11	10 occurs twice, at positions different modulo 2
	00 01 10 11	
	00 01 10 11	11 occurs twice, at positions different modulo 2

# Observation

Consider the concatenation in lexicographic order of all blocks of length  $n$ . Viewed circularly, each block of length  $n$  occurs exactly  $n$  times at positions different modulo  $n$ .

$n = 3$	000 001 010 011 100 101 110 111	000 occurs three times,
	000 001 010 011 100 101 110 111	at positions different modulo 3
	000 001 010 011 100 101 110 111	
	000 001 010 011 100 101 110 111	001 occurs three times
	000 001 010 011 100 101 110 111	at positions different modulo 3
	000 001 010 011 100 101 110 111	
	...	
	⋮	

Neither Barbier (1887) nor Champernowne (1933) noticed this.



# Observation

Not every permutation of the blocks of length  $n$  has the property,

00 10 11 01

# Perfect necklaces

**Definition** (Alvarez, Becher, Ferrari and Yuhjtman 2016)

A necklace over a  $b$ -alphabet is  $(n, k)$ -perfect if each block of length  $n$  occurs  $k$  many times at positions different modulo  $k$ , for any convention of the starting point.

The  $(n, k)$ -perfect necklaces have length  $kb^n$ .

De Bruijn sequences are exactly the  $(n, 1)$ -perfect necklaces.

## Arithmetic progressions yield perfect necklaces

Identify the blocks of length  $n$  over the  $b$ -alphabet with the set of non-negative integers modulo  $b^n$  according to representation in base  $b$ .

**Theorem** (Alvarez, Becher, Ferrari and Yuhjtman 2016)

*Let  $r$  coprime with  $b$ . The concatenation of blocks corresponding to the arithmetic sequence  $0, r, 2r, \dots, (b^n - 1)r$  yields an  $(n, n)$ -perfect necklace.*

With  $r = 1$  we have the lexicographically ordered  $(n, n)$ -perfect necklace.

## Astute graphs

Fix  $b$ -alphabet  $A$ .

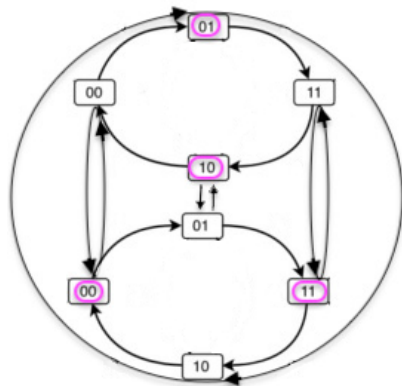
The **astute graph**  $G(b, n, k)$  is directed, with  $kb^n$  vertices.

Vertices:  $A^n \times \{0, \dots, k-1\}$

Edges:  $(w, m) \rightarrow (w', m')$  if

$$w(2..n) = w'(1..n-1) \text{ and } m' = (m + 1) \pmod k$$

This is  $G(2, 2, 2)$



## Astute graphs

Each Hamiltonian cycle in  $G(b, n, k)$  gives one  $(n, k)$ -perfect necklace.

The line graph of  $G(b, n - 1, k)$  is  $G(b, n, k)$ .

Each Eulerian cycle in  $G(b, n - 1, k)$  gives one  $(n, k)$ -perfect necklace.

Each  $(n, k)$ -perfect necklace possibly gives various Eulerian cycles in  $G(b, n - 1, k)$ .

# Normal sequences and perfect necklaces

**Theorem** (proved first by Ugalde 2000 for de Bruijn)

*The concatenation of  $(n, k)$ -perfect necklaces over a  $b$ -alphabet, for  $(n, k)$  linearly increasing, is normal to the  $b$ -alphabet.*



## A seemingly easier criterion for normality

### Theorem (Piatetski-Shapiro 1951)

Let  $x$  be an infinite sequence in a  $b$ -symbol alphabet.

If there is a constant  $C$  such that for all words  $w$ ,

$$\limsup_{n \rightarrow \infty} \frac{\# \text{occurrences of } w \text{ in } x[1, n]}{n} < Cb^{-|w|}$$

then  $x$  is normal.

# First result insertion

## Theorem (Becher 2022)

Let  $x$  be a concatenation of  $(n + 1, n)$ -perfect necklaces over  $b$ -alphabet,  $n \geq 1$ .  
There is computable sequence  $y$  normal to  $(b + 1)$ -alphabet such that  $y$  is the concatenation of  $(n + 1, n)$ -perfect necklaces over  $(b + 1)$ -alphabet for  $n \geq 1$ , and  $x$  is a subsequence of  $y$ .

Moreover, for every integer  $N$  greater than  $b$ , in between the occurrences of the new symbol in  $y$  just before and just after position  $N$  there are at most  $2b + \log_{b+1}(N)$  symbols.





# Insertion in perfect necklaces

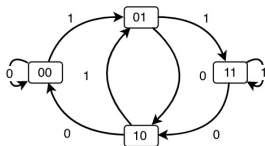
**Theorem** (follows from Becher and Cortés 2020)

*For every  $(n + 1, n)$ -perfect necklace  $v$  over a  $b$ -alphabet there is an  $(n + 1, n)$ -perfect necklace  $w$  over  $(b + 1)$ -alphabet such that  $v$  is a subsequence of  $w$ .*

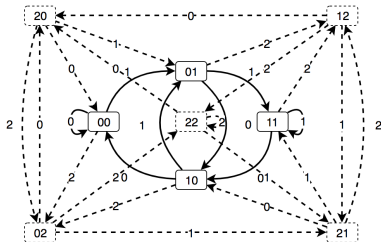
*Moreover, for each such  $v$  there is  $w$  satisfying that for any  $n + 2b - 1$  consecutive symbols there is at least one occurrence of the new symbol.*

# Proof of theorem on insertion in perfect necklaces

Given  $(n + 1, n)$ -perfect necklace over  $b$ -alphabet, pick an Eulerian cycle in  $G(b, n, n)$  corresponding to it.

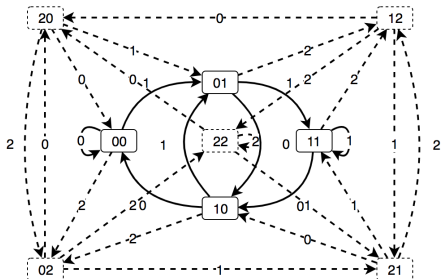


$G(b, n, n)$  is a subgraph of  $G(b + 1, n, n)$



Pictures here are  $G(b, n, 1)$  and  $G(b + 1, n, 1)$  instead of  $G(b, n, n)$  and  $G(b + 1, n, n)$

Consider  $G(b + 1, n, n) \setminus G(b, n, n)$  (the dotted lines)



Every Eulerian cycle is the union of disjoint cycles.

**Without the small gap condition:**

just use Euler-Hierholzer's algorithm for joining cycles

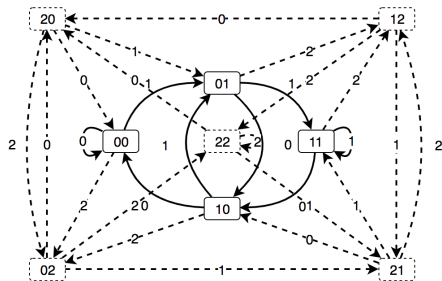
**With the small gap condition** is more delicate.

The resulting Eulerian cycle in  $G(b + 1, n, n)$  is an  $(n + 1, n)$ -perfect necklace on  $(b + 1)$ -alphabet with the wanted properties.

## Augmenting graph

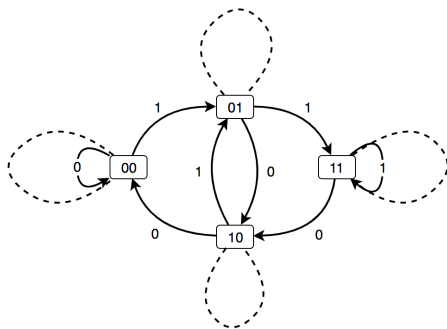
The **augmenting graph**  $A(b+1, n, n)$  has exactly all the vertices of  $G(b+1, n, n)$  and all the edges of  $G(b+1, n, n) \setminus G(b, n, n)$ .

Picture for  $k = 1$ , de Bruijn case.



# Petals

Partition the augmenting graph in  $nb^n$  disjoint cycles, called **petals**, such that in each of them there is one vertex in  $G(b, n, n)$ .



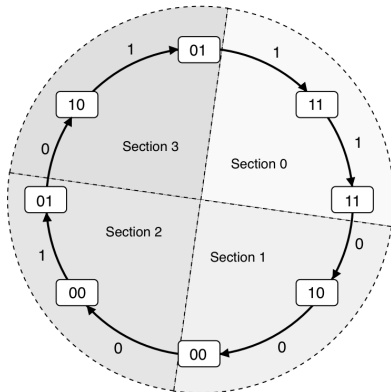
## The small gap condition

Consider an  $(n + 1, n)$ -perfect necklace and a starting position.

Pick corresponding Eulerian cycle in  $G(b, n, n)$ , with edges  $e_1, \dots, e_{nb^{n+1}}$ .

Divide it in  $nb^n$  consecutive sections, each consisting of  $b$  edges.

Identify each section with the  $b$  target vertices in it.



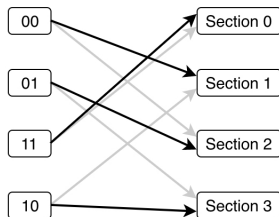
## A matching problem

The astute graph  $G(b, n, n)$  has  $nb^n$  vertices.

An Eulerian cycle in  $G(b, n, n)$  has  $nb^n$  sections with  $b$  vertices each.

We need to choose **one vertex in each section**, and they are all different.

We pose it as a **matching** problem.



By Hall's marriage theorem there is a **perfect matching**.

Find it computing the maximum flow in a network, Edmonds-Karp algorithm  $\square$

# Any symbol insertion in constructed normal sequences

We have just showed:

Starting from some particular constructed normal sequences, we can insert symbols and obtain a normal sequence in the enlarged alphabet, with the small gap condition.



## Second result on insertion

### Theorem (Zylber 2017)

For every sequence  $x$  normal to  $b$ -alphabet there exists a sequence  $y$  normal to  $(b + 1)$ -alphabet such that  $\text{retract}(y) = x$ , where  $\text{retract}(\cdot)$  removes the occurrences of the symbol ' $b$ '.



In general, the sequence  $y$  is **not** computable from  $x$ , but from an upper bound of the speed of convergence to normality of  $x$ .

## Proof idea

Given a sequence over  $b$ -alphabet, insert ' $b$ ' in the positions prescribed by the concatenation of lexicographically ordered  $(n, n)$ -perfect necklaces over  $(b + 1)$ -alphabet

perfect	00	01	02	10	11	12	20	21	22
wildcards	**	**	*2	**	**	*2	2*	2*	22

# Proof idea

Given a sequence over  $b$ -alphabet, insert ' $b$ ' in the positions prescribed by the concatenation of lexicographically ordered  $(n, n)$ -perfect necklaces over  $(b + 1)$ -alphabet

perfect	00	01	02	10	11	12	20	21	22
wildcards	**	**	*2	**	**	*2	2*	2*	22
arbitrary	$a_1a_2$	$a_3a_4$	$a_5a_6$	$a_7a_8$	$a_9a_{10}$	$a_{11}a_{12}$			
insertion	$a_1a_2$	$a_3a_4$	$a_5$ 2	$a_6a_7$	$a_8a_9$	$a_{10}$ 2	2 $a_{11}$	2 $a_{12}$	22

## Discrete discrepancy on aligned occurrences

$$\Delta_{b,\ell}(u) = \max_{\text{block } v \text{ of length } \ell} \left| \frac{\text{number of aligned occurrence of } v \text{ in } u}{\lfloor |u|/\ell \rfloor} - \frac{1}{b^\ell} \right|.$$

### Fact

If  $p$  divides  $P$  and  $\Delta_{b,P} < \epsilon$  then  $\Delta_{b,p} < \text{bound}(b, \epsilon, p, P)$ .

Proved by Pillai (1940):

A sequence  $a_1 a_2 \dots$  is normal to  $b$ -alphabet if for every length  $\ell$

$$\lim_{n \rightarrow \infty} \Delta_{b,\ell}(a_1 \dots a_{\ell n}) = 0$$

In general, given  $a_1 a_2 \dots$  the function  $f : \mathbb{Q} \rightarrow \mathbb{N}$  such that  $f(\epsilon) = n$  when  $n$  is the minimum such that  $\Delta_{b,\ell}(a_1 \dots a_n) < \epsilon$  is **not** computable.

# Insertion and discrete discrepancy

Let  $\ell_n$  be the number of non- $b$  symbols in  $(n, n)$ -perfect necklace on  $(b + 1)$ -alphabet.

Let  $e_n(u)$  be the sequence that results from inserting  $b$ 's in  $u$  according to the ordered  $(n, n)$ -perfect necklace on  $(b + 1)$ -alphabet.

## Lemma (crucial)

*For every  $n$ , there is  $c_n$  such that for every finite  $u$  over  $b$ -alphabet*

$$\begin{aligned} \text{if } \Delta_{b, \ell_n}(u) < \epsilon \\ \text{then } \Delta_{b+1, n}(e_n(u)) < c_n \epsilon \end{aligned}$$

# Zylber's construction

Given sequence  $x$  normal to  $b$ -alphabet,

- Partition  $x$  in  $u_1, u_2, u_3, \dots$  where  $|u_n| = t_n \ell_{2^n}$  and  $t_n$  is large enough so that  $\Delta_{b, \ell_{2^n}}(u_n)$  is small.

Recall  $\ell_{2^n}$  is the number of non- $b$ 's in  $(2^n, 2^n)$ -perfect necklace on  $(b+1)$ -alphabet.

- By the crucial Lemma  $\Delta_{b+1, 2^n}(e_{2^n}(u_n))$  is small.
- Discrete discrepancy  $\Delta_{b+1, P}$  also controls  $\Delta_{b+1, p}$  when  $p$  divides  $P$ .

The wanted sequence is  $y = e_{2^1}(u_1)e_{2^2}(u_2)e_{2^3}(u_3)\dots$

By Piatetski-Shapiro theorem,  $y$  is normal to  $(b+1)$ -alphabet. □

## Open problems

- ▶ What are other forms of insertion transferring normality from  $b$ -alphabet to  $(b + 1)$ -alphabet?

- ▶ Discrepancy analysis for different forms of insertion

Compare discrepancy  $(b^n x \bmod 1)_{n \geq 0}$  and  $((b + 1)^n y \bmod 1)_{n \geq 0}$  where  $y$  results from insertion in  $x$ .

Fukuyama and Hiroshima in 2012 gave metric discrepancy results for subsequences of  $(b^n x \bmod 1)_{n \geq 0}$ ,

- ▶ How can we change effective Hausdorff dimension while enlarging the alphabet?

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